



# CS F211: DATA STRUCTURES & ALGORITHMS (2<sup>ND</sup> SEMESTER 2024-25) TREE ADT

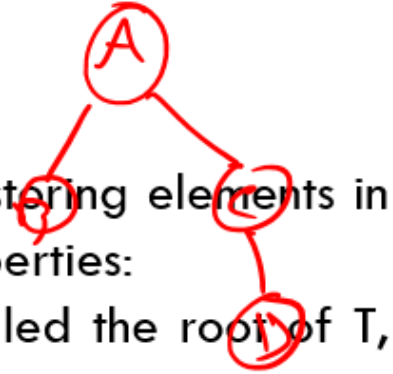
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# TREES: NON-LINEAR DATA STRUCTURES

- In computer science, what is a tree?

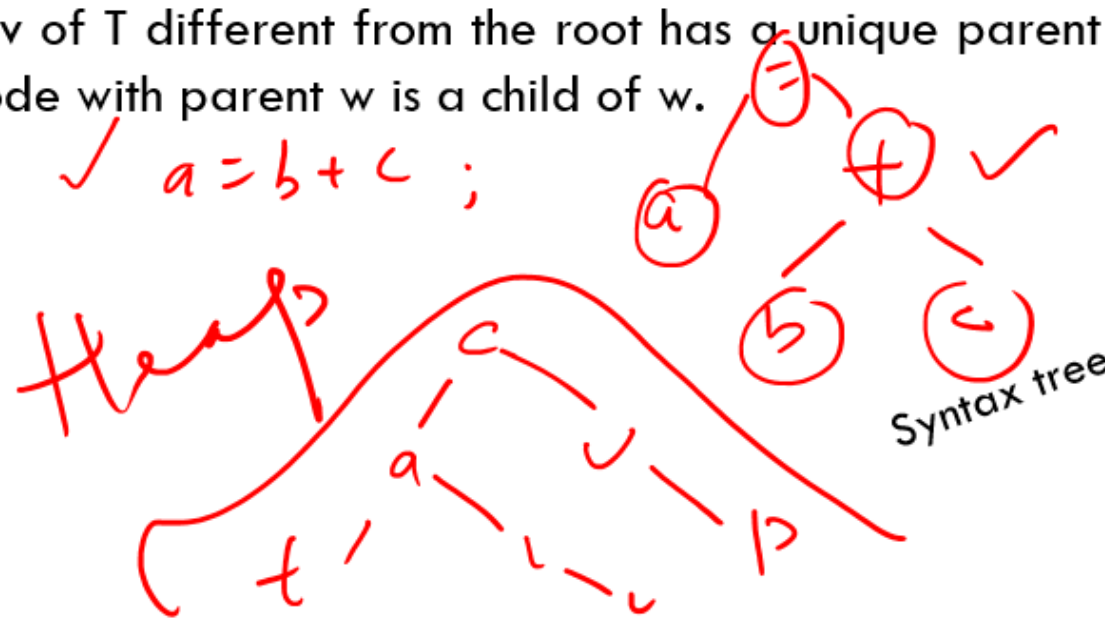
Formally, we define tree  $T$  to be a set of nodes storing elements in a **parent-child relationship** with the following properties:

- If  $T$  is nonempty, it has a special node, called the **root** of  $T$ , that has no parent.
- Each node  $v$  of  $T$  different from the root has a unique parent node  $w$ ; every node with parent  $w$  is a child of  $w$ .



Applications?

Tree



# TREE TERMINOLOGIES AND PROPERTIES

Root: ???

Internal node: ???

External node: ???

Ancestors of a node: ???

Level of a node or Depth of a node:  
???

Height of a node and height of tree:  
???

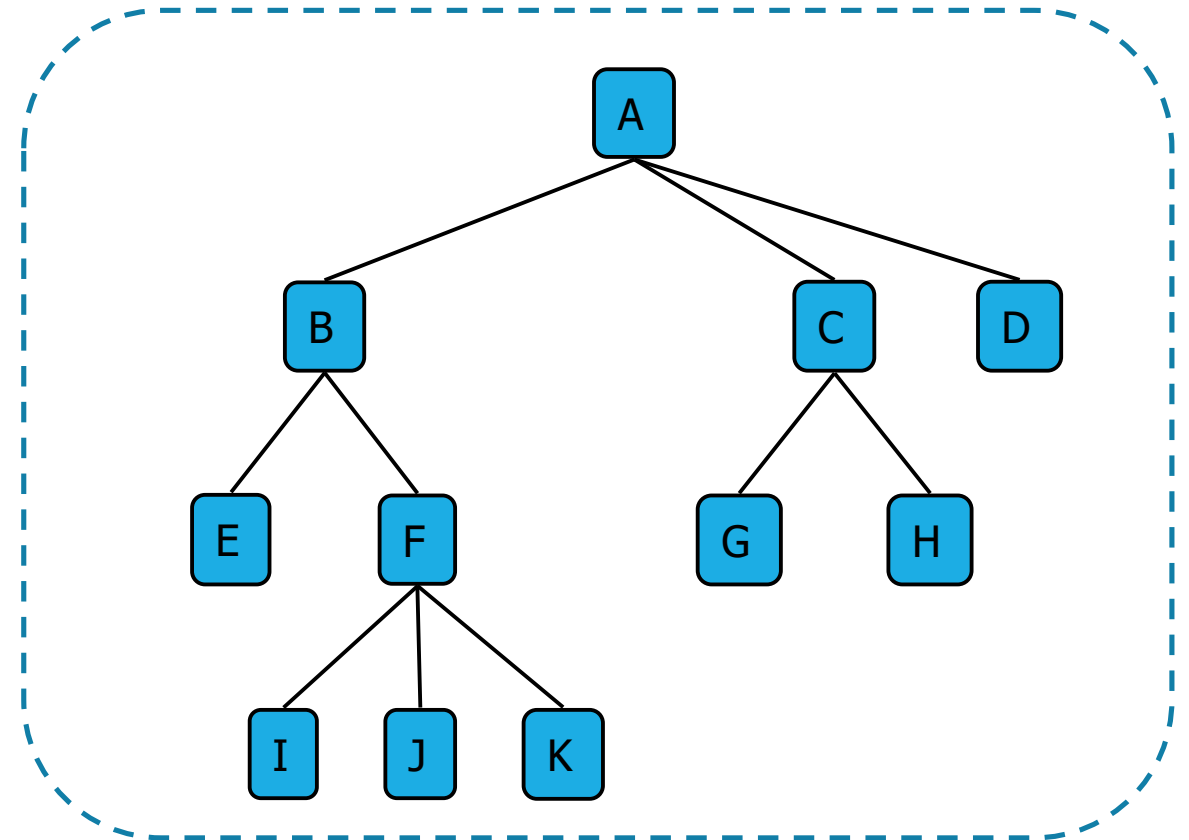
Descendants of a node: ???

What is a sub-tree?

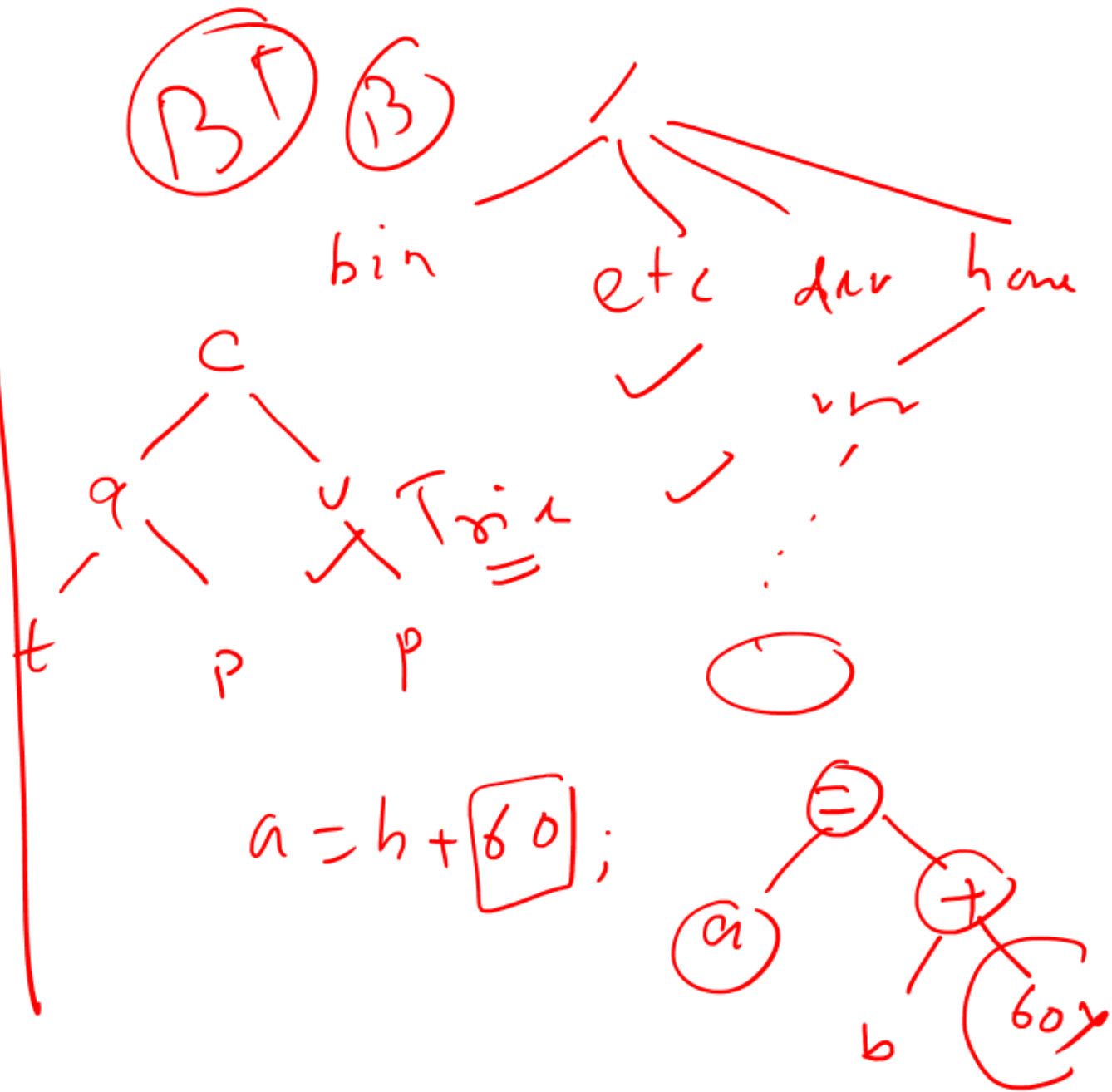
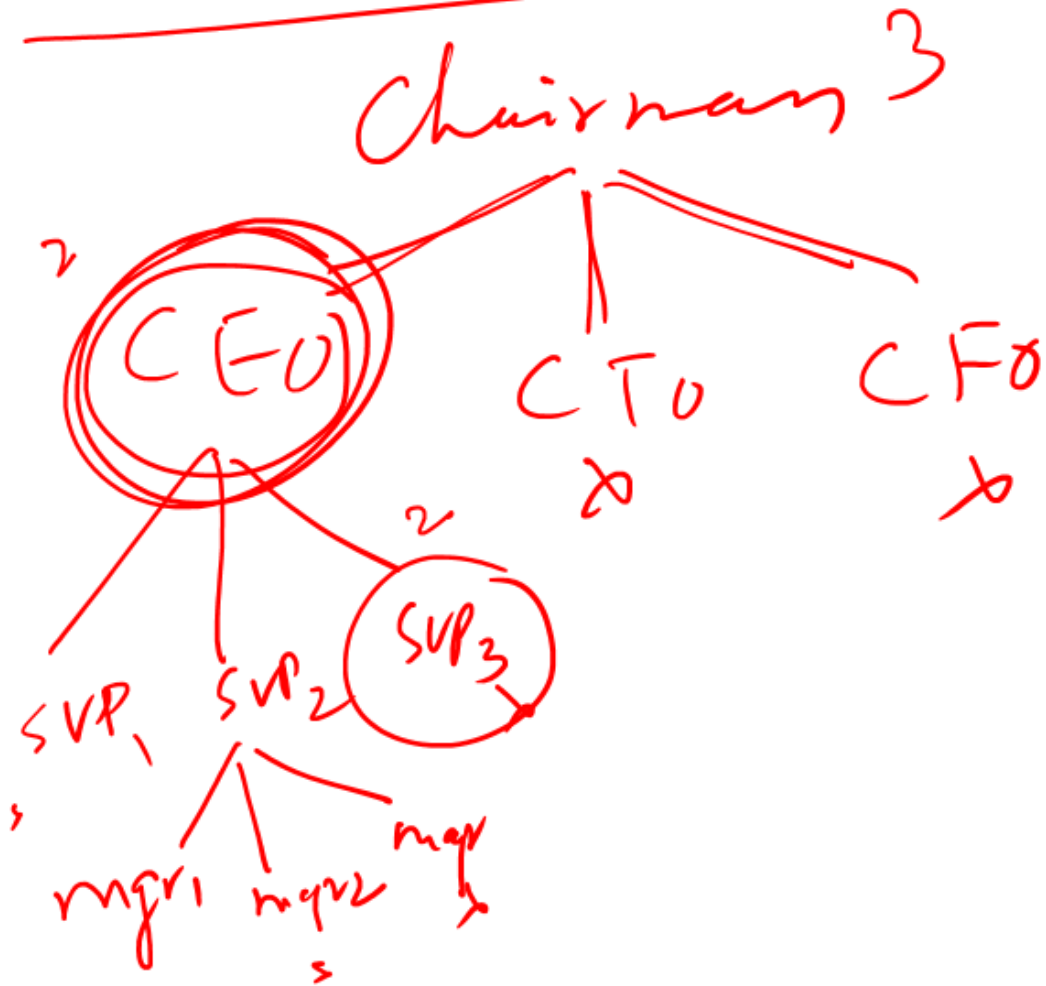
No of edges???

Degree of a node???

Forest???

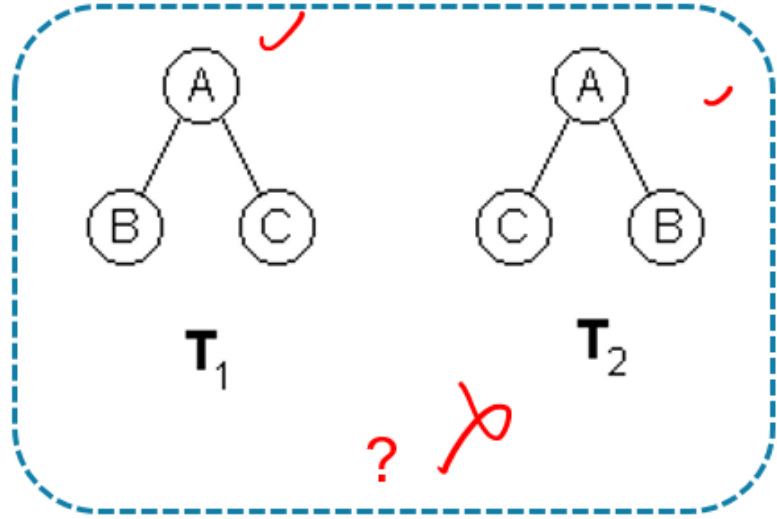
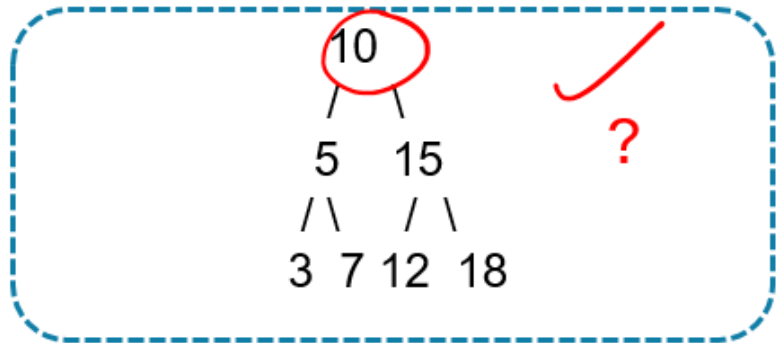
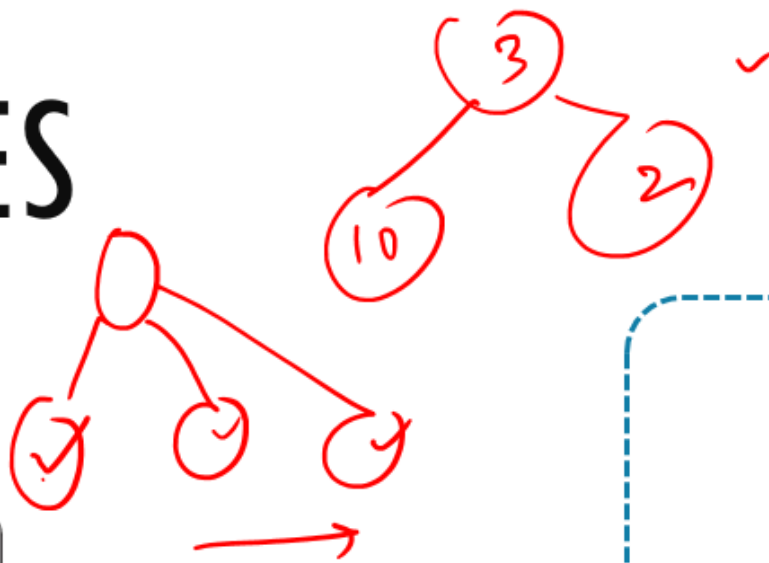
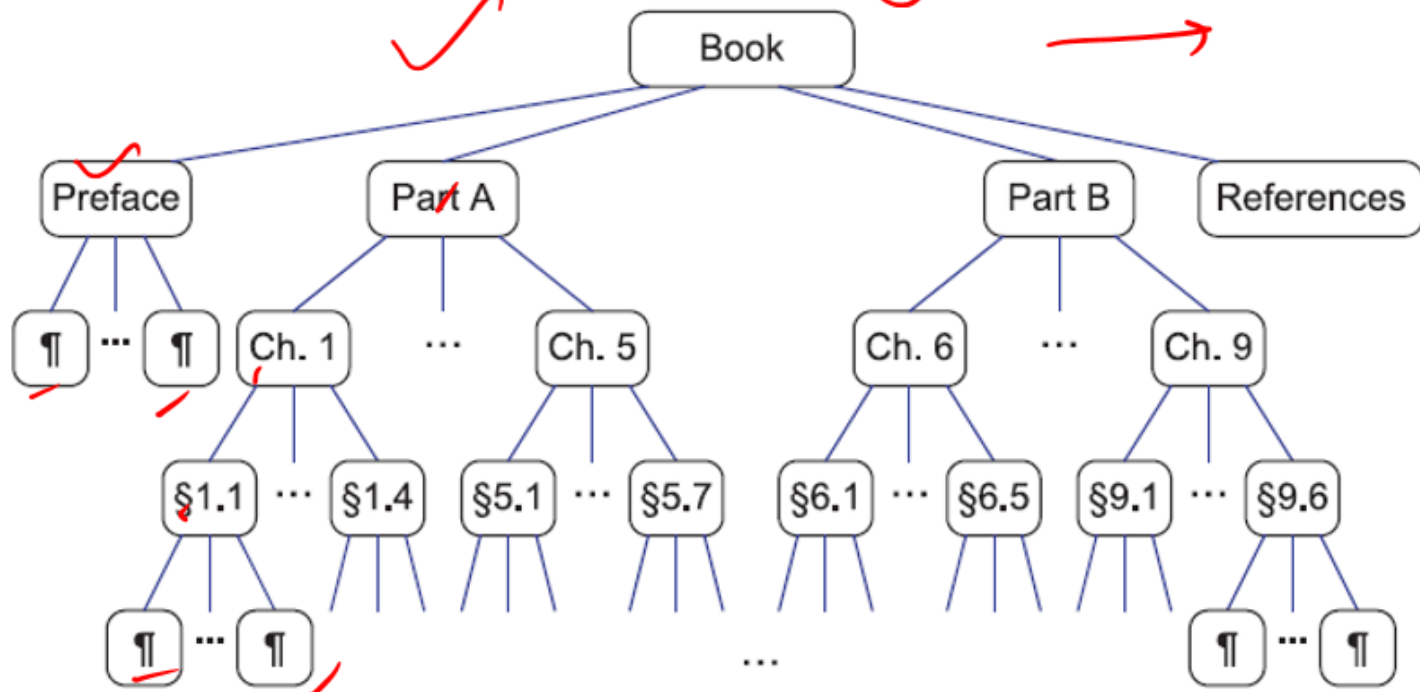


# TRY YOURSELF...

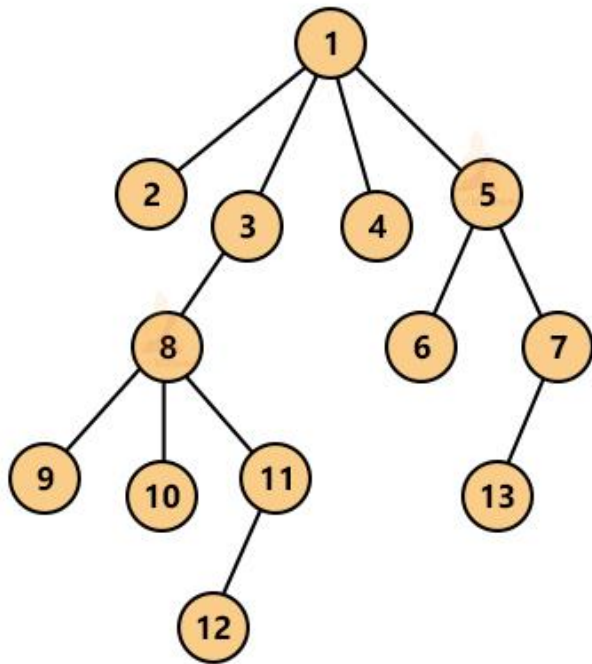


# ORDERED TREES

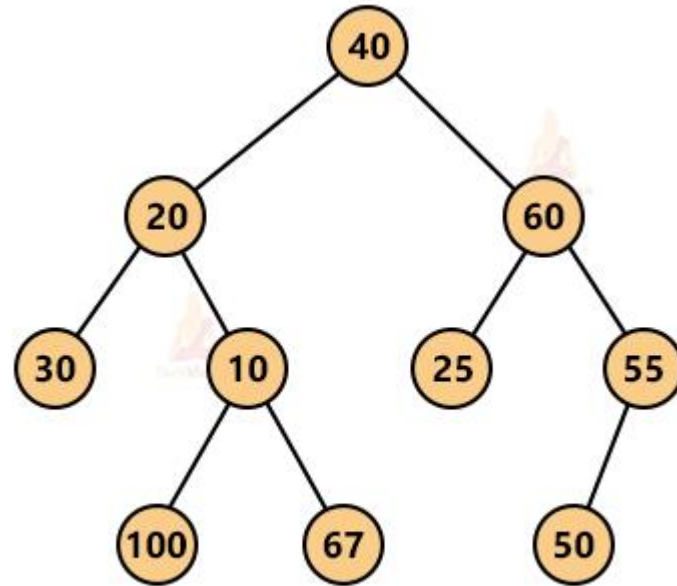
What are Ordered Trees?



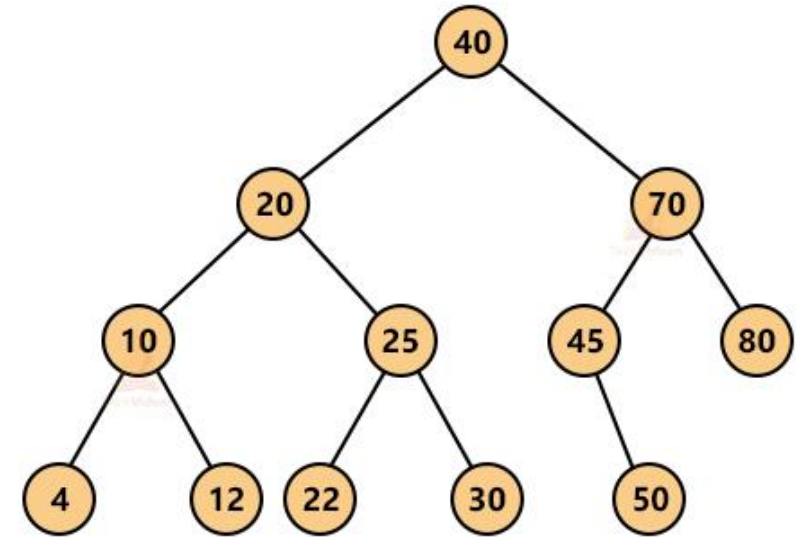
# TYPES OF TREES IN DATA STRUCTURES



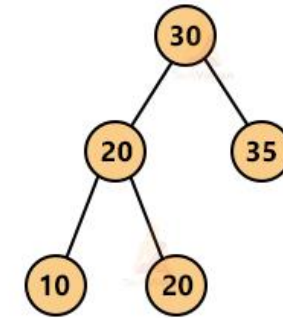
(A general tree)



(A binary tree)



(A binary search tree)



(An AVL trees)

# TREE ABSTRACT DATA TYPE (ADT)

- Generic methods:

- integer `size()`
- boolean `empty()`

- Accessor methods:

- position `root()`
- list<position> `positions()`

- Position-based methods:

- position `p.parent()`
- list<position> `p.children()`

- Query methods:

- boolean `p.isRoot()`, boolean `p.isExternal()`

- Additional update methods may be defined by data structures implementing the Tree ADT.

```
template <typename E>
class Position<E> {
public:
    E& operator*();
    Position parent() const;
    PositionList children() const;
    bool isRoot() const;
    bool isExternal() const;
};
```

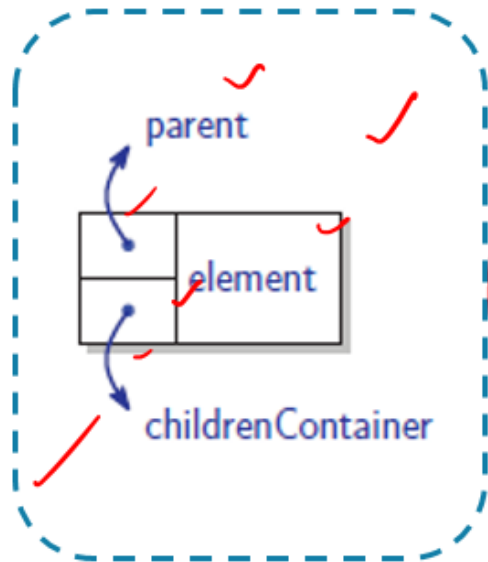
(An informal interface  
for a position in a  
tree)

```
template <typename E>
class Tree<E> {
public:
    class Position;
    class PositionList;
public:
    int size() const;
    bool empty() const;
    Position root() const;
    PositionList positions() const;
};
```

// base element type  
// public types  
// a node position  
// a list of positions  
// public functions  
// number of nodes  
// is tree empty?  
// get the root  
// get positions of all nodes

(An informal interface for the tree ADT)

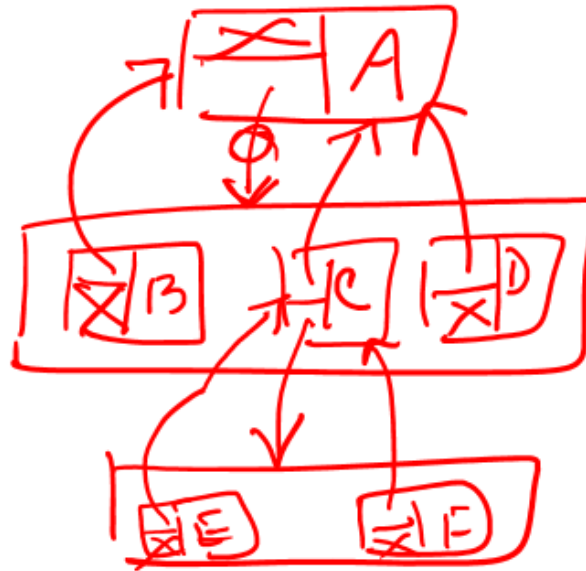
# A LINKED STRUCTURE FOR GENERAL TREES



(The node structure:  
An example)



(The portion of the data structure associated  
with root node and its children)



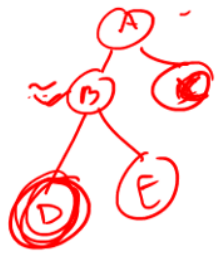
Operation	Time
isRoot, isExternal	$O(1)$
parent	$O(1)$
children( $p$ )	$O(c_p)$
size, empty	$O(1)$
root	$O(1)$
positions	$O(n)$

(Running times of linked  
tree structure)



# DEPTH AND HEIGHT OF NODE IN A TREE

if  $n \neq 0$



Algorithm `depth(T, p):`

if `p.isRoot()` then  
return 0;

else

return  $1 + \text{depth}(T, p.\text{parent}())$ ;

Algorithm `height(T):`

$h = 0$ ;

for each  $p \in T.\text{positions}()$  do

if `p.isExternal()` then

$h = \max(h, \text{depth}(T, p))$

return  $h$

$$O\left(n + \sum_{p \in E} (1 + d_p)\right)$$

$$\rightarrow O\left(n + \sum_{p \in E} (1) + \sum_{p \in E} (d_p)\right)$$

$$\begin{array}{ccc} \downarrow & & \downarrow \\ n & & 0+1+2+\dots+(n-1) \end{array}$$

$$\frac{n(n-1)}{2}$$

$\rightarrow O(n^2)$

Algorithm `height(T, p):`

if `p.isExternal()` then

return 0;

else

$h = 0$ ;

for each  $q \in p.\text{children}()$  do

$h = \max(h, \text{height}(T, q))$

return  $1+h$ ;

The recursive function spends

$O(1 + c_p)$  time at each node  $p$ .

$$O\left(\sum_p (1 + c_p)\right)$$

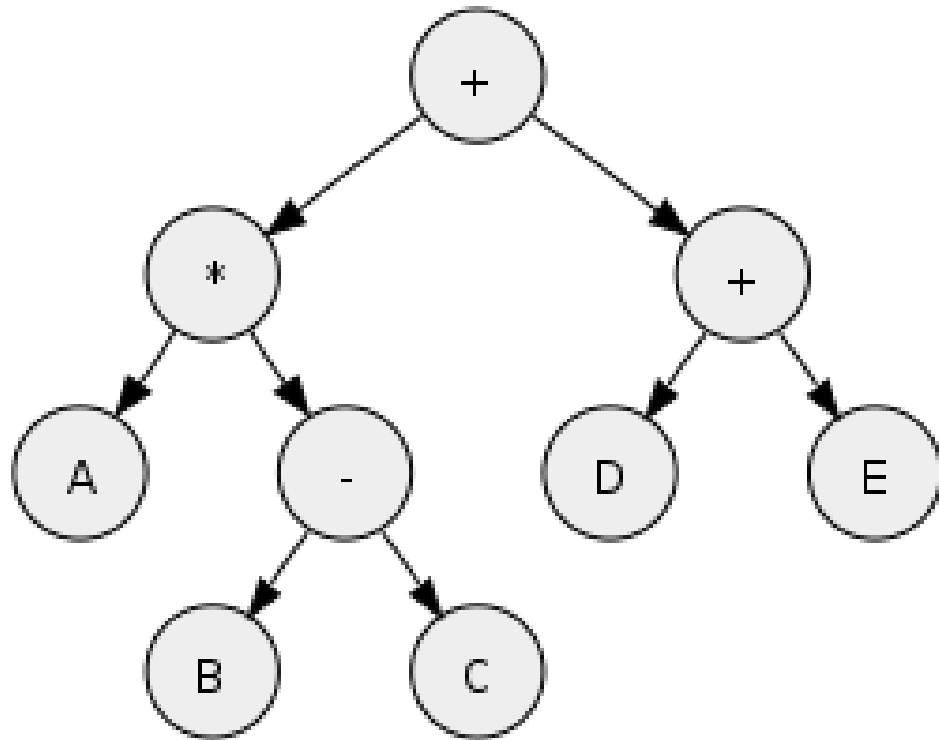
$$\begin{array}{cc} \downarrow & \downarrow \\ \sum_p (1) & \sum_p (c_p) \\ n & n-1 \end{array}$$

$\rightarrow O(n)$

Run-time Complexity?

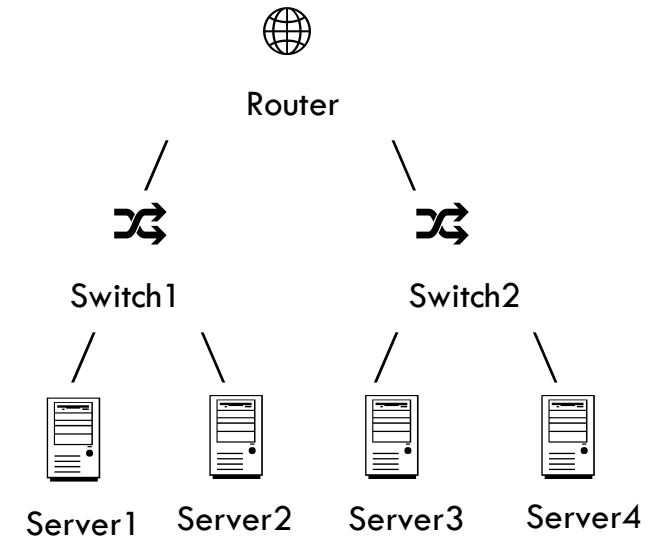
- $O(d_p)$  where  $d_p$  is the depth of node  $p$ .
- In worst case,  $d_p$  can be 'n' and hence worst case complexity is  $O(n)$ .

# TREE TRAVERSAL ALGORITHMS



(Expression Tree)

Pre-order: + \* A - B C + D E    Post-order: A B C - \* D E + +

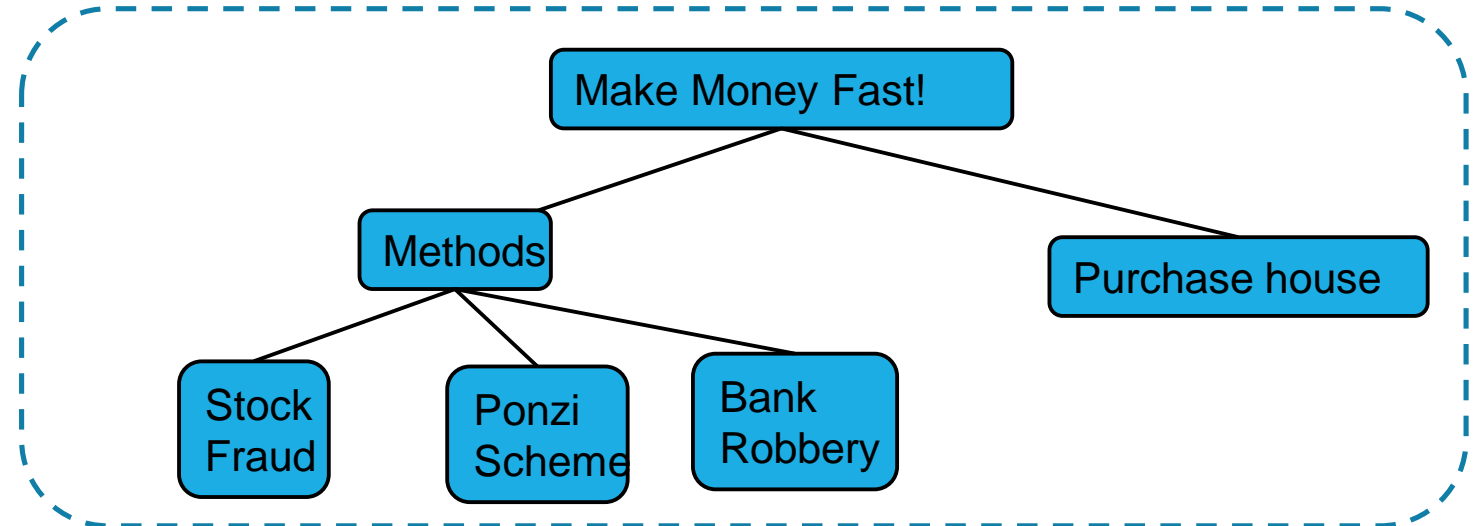


(Network shutdown using **what-order** traversal?)

Post order: Server1, Server2, Switch1, Server3, Server4, Switch2, Router.

# TREE TRAVERSAL: DEPTH-FIRST (PRE-ORDER)

- In a preorder traversal, a node is visited **before** its descendants.
- Applications: print a structured document, get the prefix expression on an expression tree.



```
20 void printPreorder(struct Node* node)
21 {
22     if (node == NULL)
23         return;
24     cout << node->data << " ";
25
26     printPreorder(node->left);
27
28     printPreorder(node->right);
29 }
```

```
Preorder traversal of binary tree is
1 2 4 5 3
```

Algorithm *preOrder* ( $T, p$ )  
{ visit ( $p$ );  
preorder( $p$ ->left);  
preorder( $p$ ->right);  
}

← Lab 8 after mid sem

```
PreOrderIterative(root){
if (root == NULL)
return;
Create stack  $S = \Phi$ ;
Push root to  $S$ ;
While Stack  $S \neq \Phi$  {
N=Pop from  $S$ ;
Print N.value;
if (N.right  $\neq$  NULL)
Push N.right;
if (N.left  $\neq$  NULL)
Push N.left;
}
```

# POST-ORDER TRAVERSAL

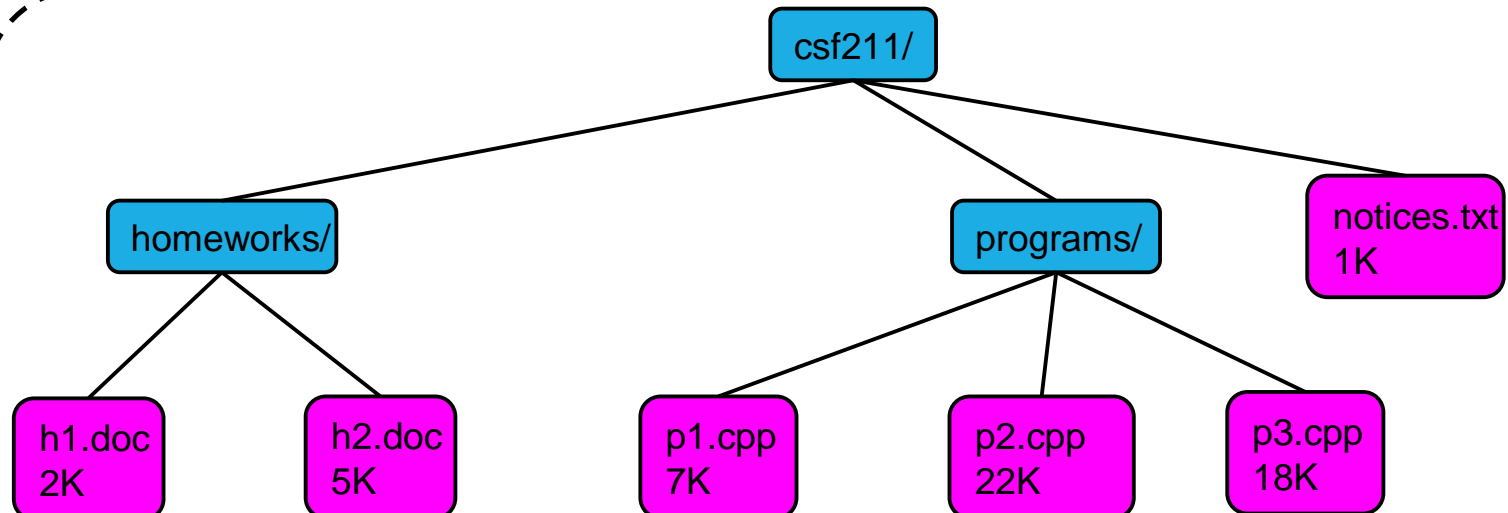
- In a postorder traversal, a node is visited **after** its descendants.
- Application: compute space used by files in a directory and its subdirectories, delete the tree, compute postfix expression...

```
23 void printPostorder(struct Node* node)
24 {
25     if (node == NULL)
26         return;
27
28     printPostorder(node->left);
29
30     printPostorder(node->right);
31
32     cout << node->data << " ";
33 }
```

Lab 8 after mid sem

```
Postorder traversal of binary tree is
4 5 2 3 1
```

Algorithm `postOrder(v){`  
for each `child w of v`  
    `postOrder (w);`  
`visit(v);`  
}



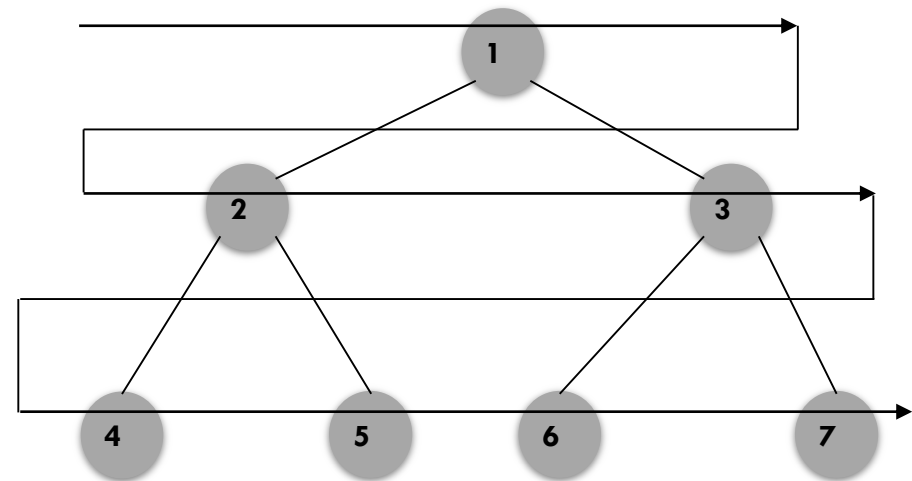
# TREE TRAVERSAL CONTINUED... LEVEL ORDER

```
LevelOrderTraversal(root) {  
  if (root == NULL) then  
    return;  
  Q ←  $\Phi$   
  Enqueue (root);  
  while (Q ≠  $\Phi$ ) {  
    node = Dequeue from Q;  
    Print node.value // Process the current node  
    if node.left ≠  $\Phi$   
      Enqueue node.left into Q;  
    if node.right ≠  $\Phi$   
      Enqueue node.right into Q;  
  }  
}
```

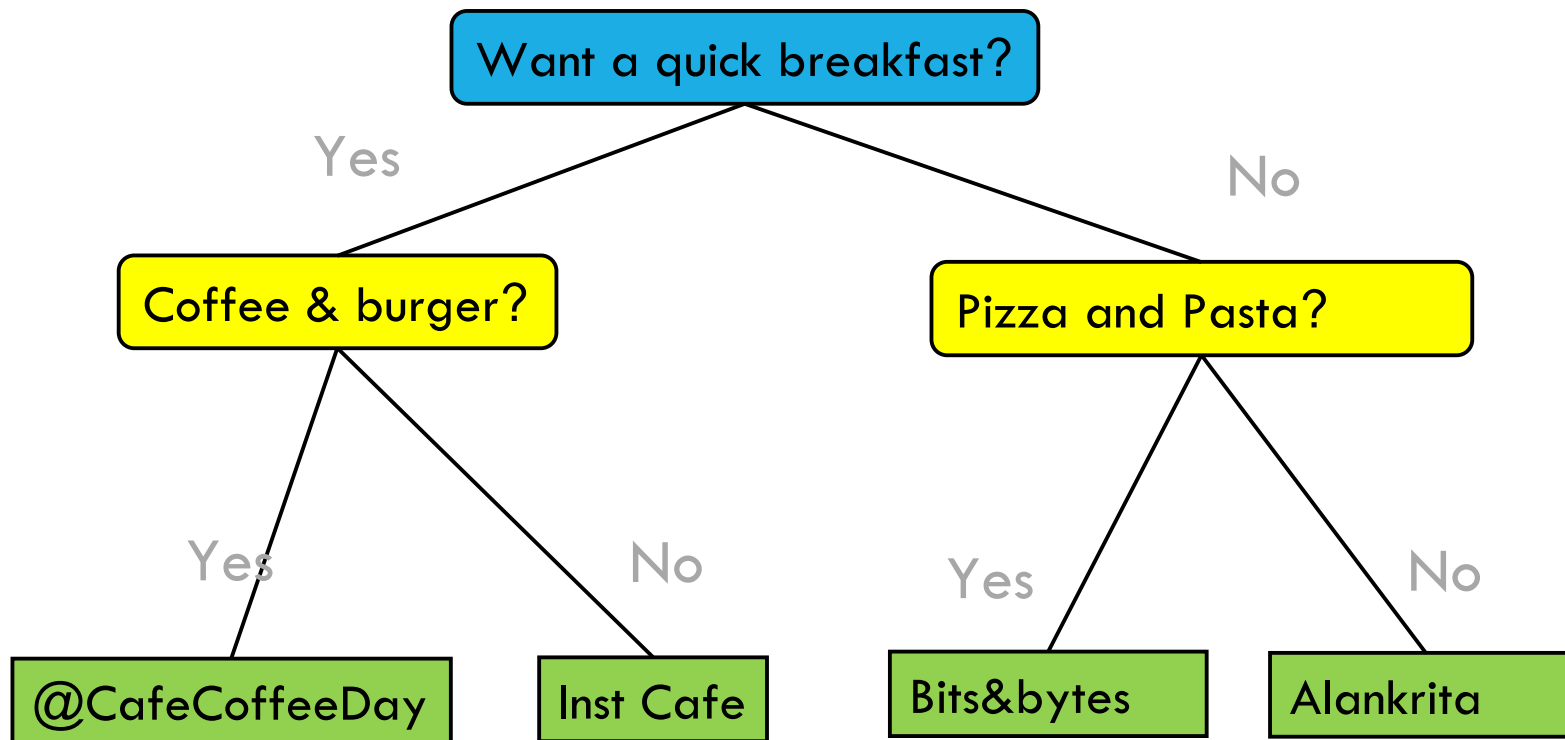
## Application:

Construct a binary tree from an array using a level-order traversal:

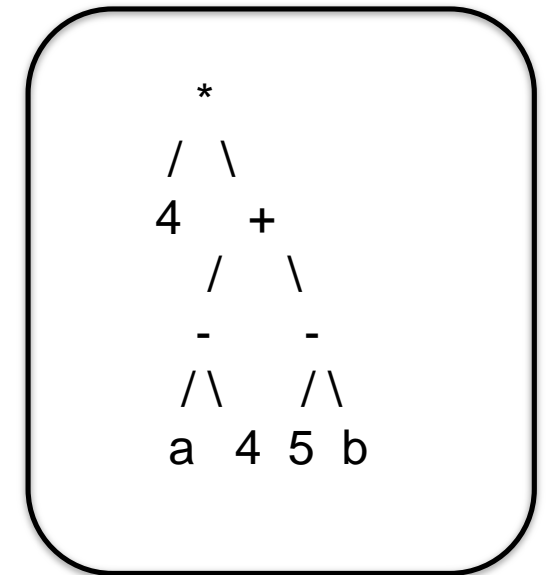
Let the array be: Arr = [1, 2, 3, 4, 5, 6, 7]



# BINARY TREES: EXAMPLE USAGES

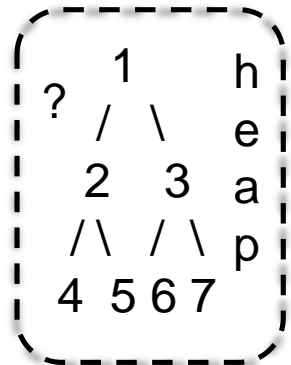
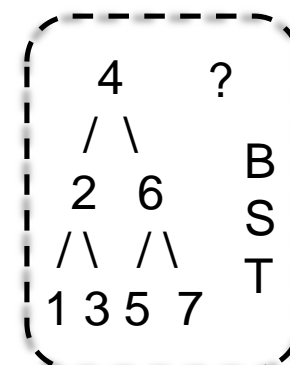
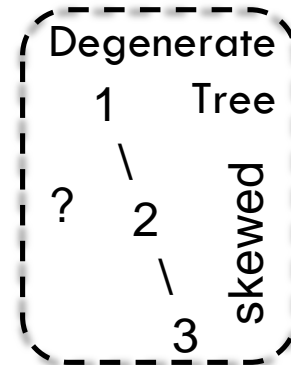
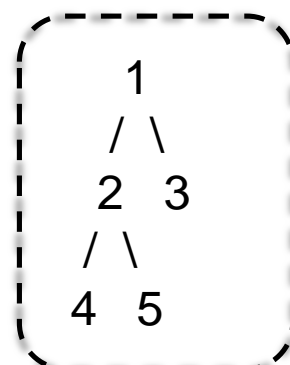
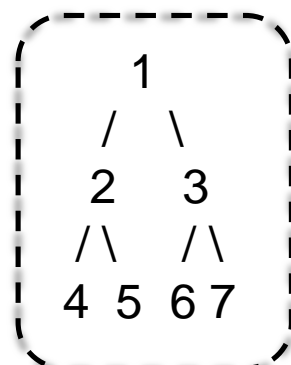
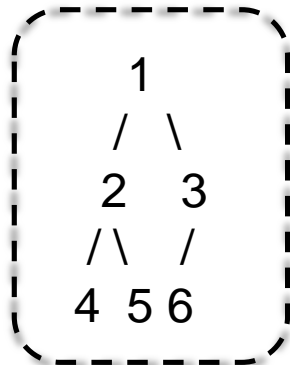
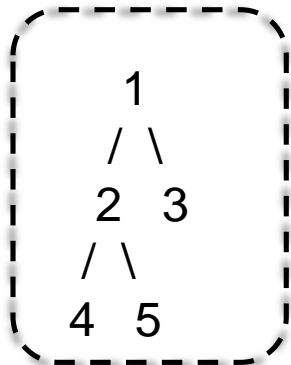
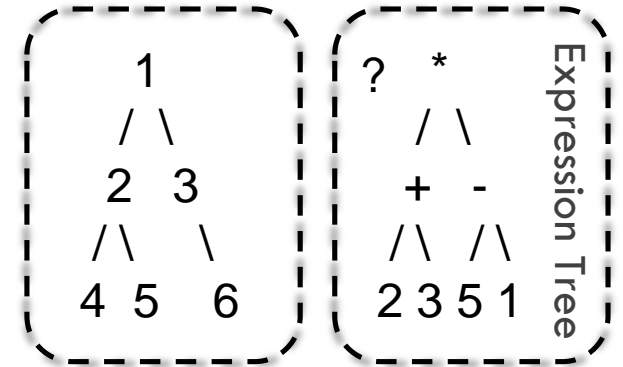


Expression Tree:  
 $4 \times ((a - 4) + (5 - b))$



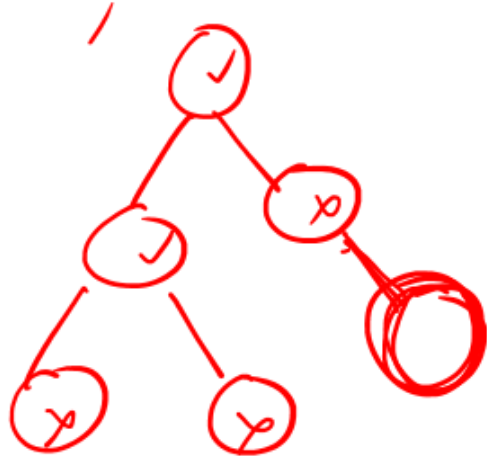
# DEFINITIONS: BINARY TREE

- A hierarchical data structure in which each node has **at most** two children.
- **Full Binary Tree**: Every node has either 0 or 2 children.
- **Complete Binary Tree**: All levels are fully filled except possibly the last level, which is filled from left to right.
- **Perfect Binary Tree**: All interior nodes have exactly two children, and all leaves are at the same depth.
- **Balanced Binary Tree**: The height of the left and right subtrees of every node differs by no more than 1.

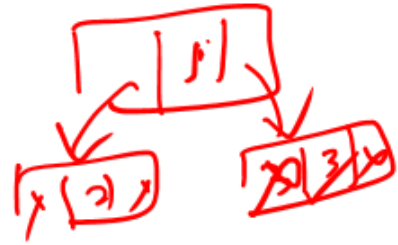


# BINARY TREES: TRY YOURSELF!

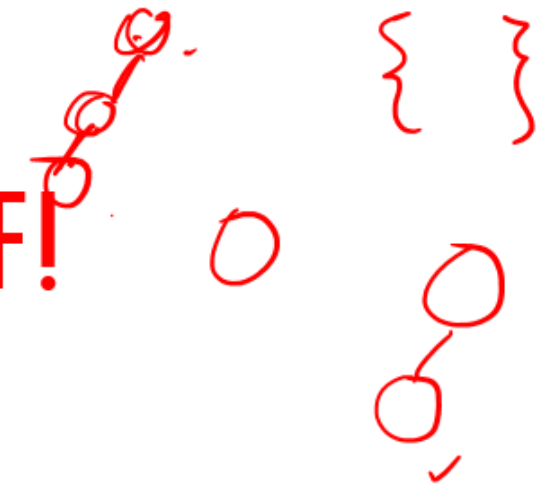
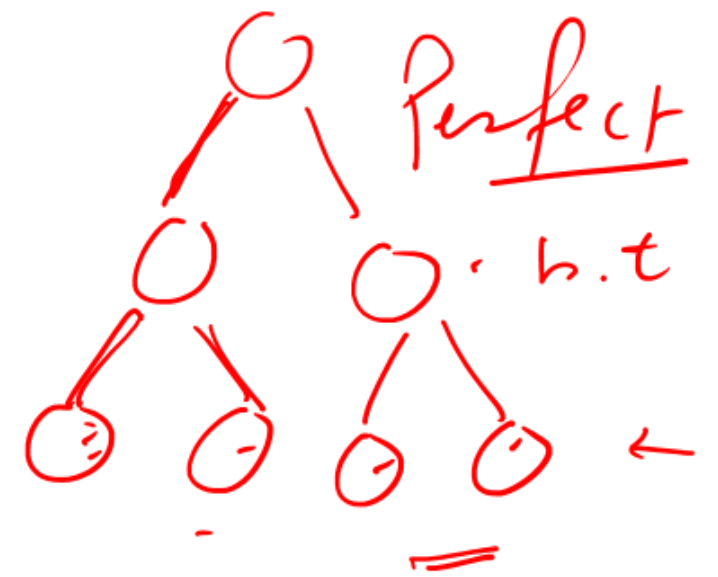
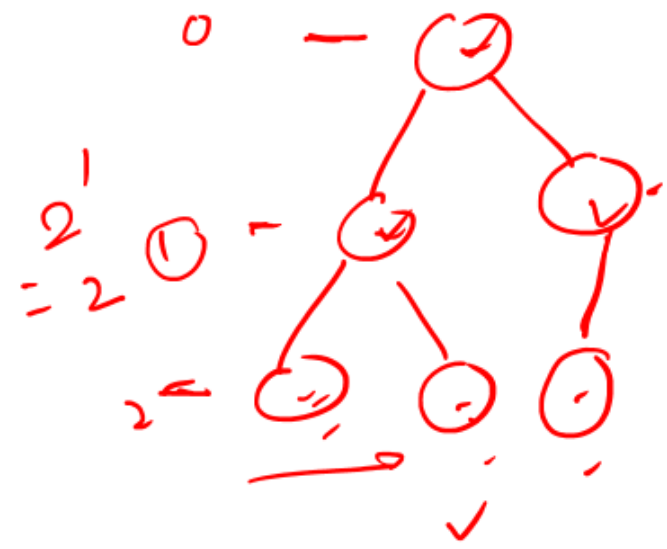
MSI ✓ full binary tree



✓ ✓ full b.t. ✗



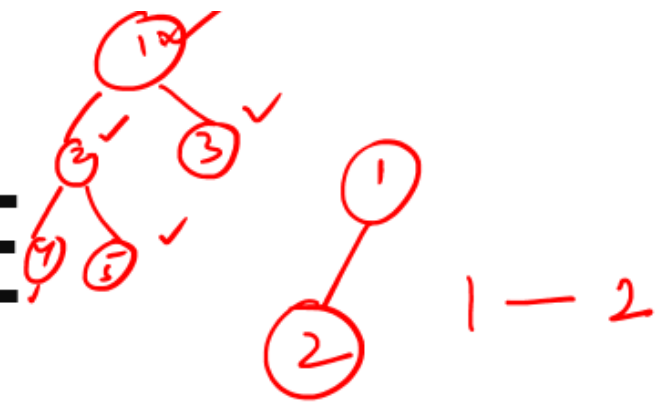
Complete b.t





# PROPERTIES OF A BINARY TREE

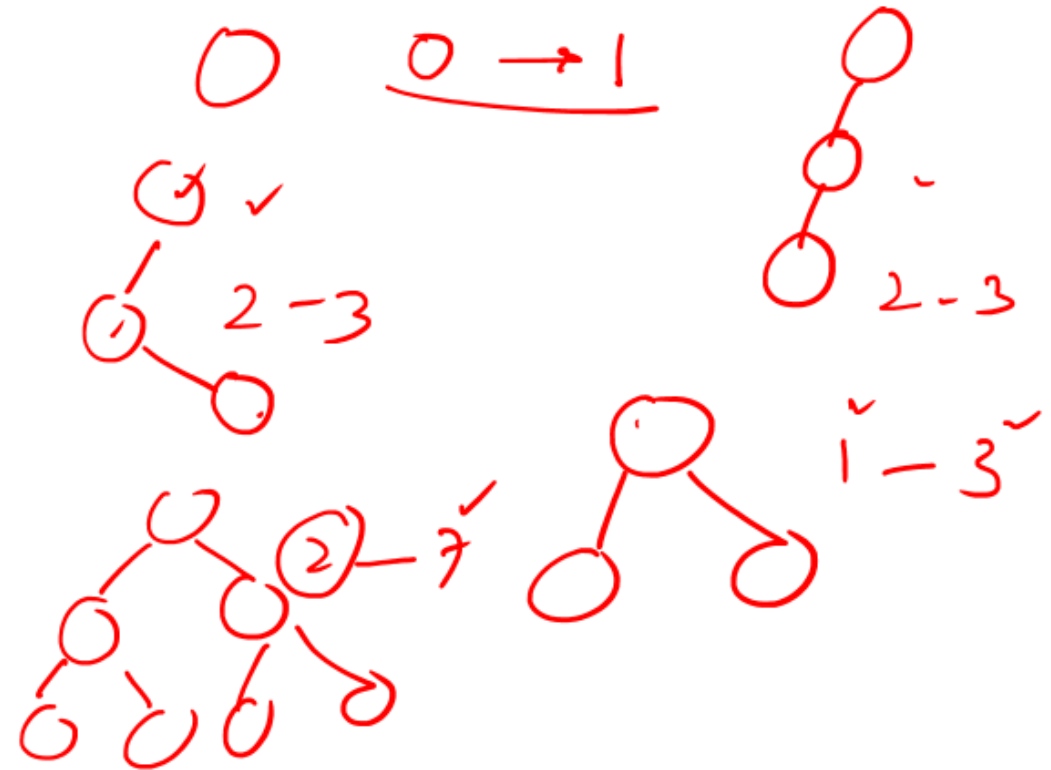
$$2 + 1 = 3$$



## Notation:

**n**: number of nodes, **e**: number of external nodes, **i**: number of internal nodes, and **h**: height

- Minimum number of nodes in a binary tree with height  $h = ?$   $h + 1$
- Maximum number of nodes in a binary tree with height  $h = ?$   $2^{h+1} - 1$
- Total number of leaf nodes in a binary tree = ? Nodes with 2 children + 1 ✓
- Maximum number of nodes at any level 'l' in a binary tree = ?  $2^l$



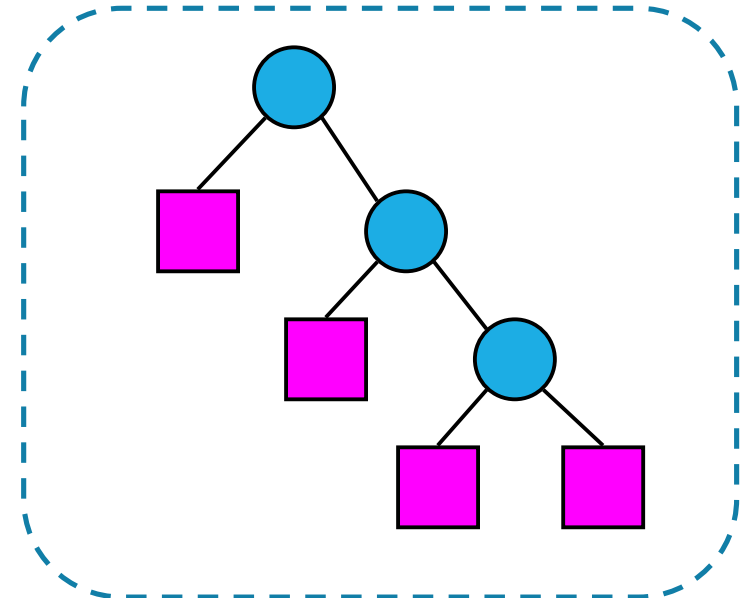
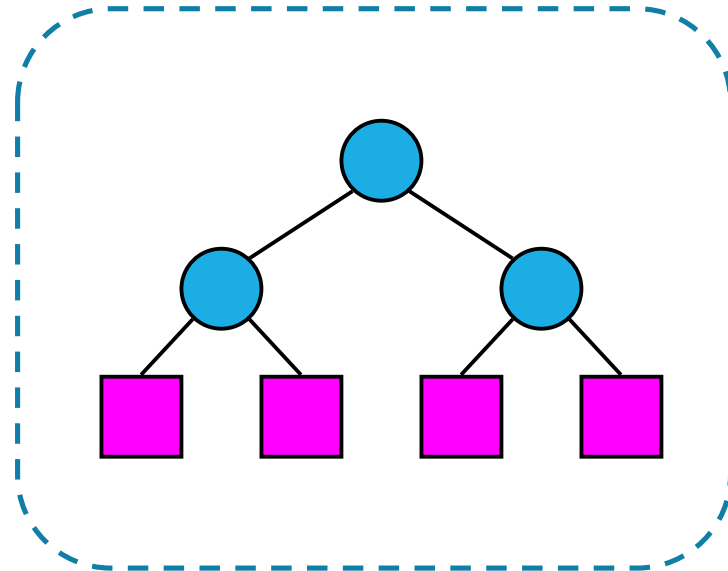
# CONTINUED...

## Notation:

$n$ : number of nodes,  $e$ : number of external nodes,  $i$ : number of internal nodes, and  $h$ : height

- Properties:

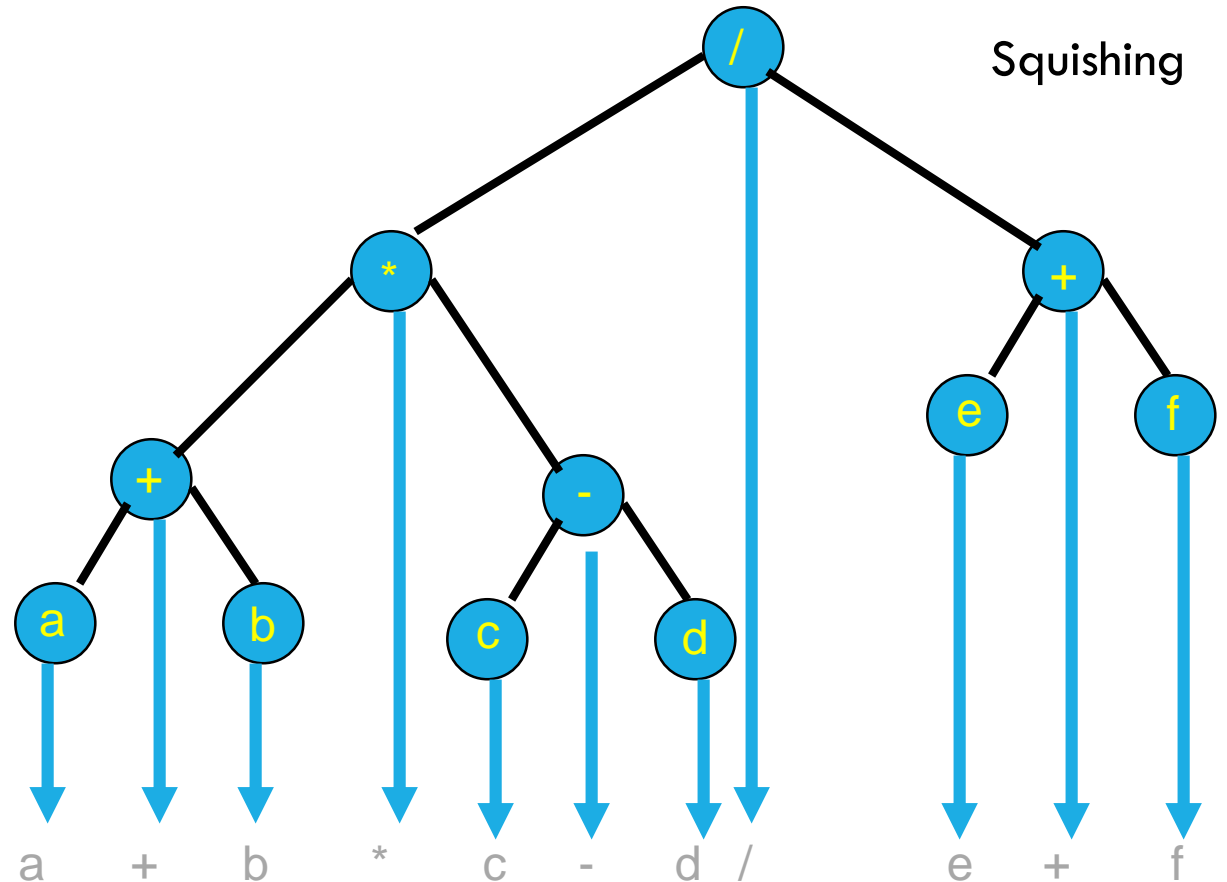
- $e = i + 1$
- $n = 2e - 1$
- $h \leq i$
- $h \leq (n - 1)/2$
- $e \leq 2^h$
- $h \geq \log_2 e$
- $h \geq \log_2 (n + 1) - 1$



# BINARY TREE TRAVERSAL: IN-ORDER

```
template <class T>
void inOrder(binaryTreeNode<T> *t)
{
    if (t != NULL)
    {
        inOrder(t->leftChild);
        visit(t);
        inOrder(t->rightChild);
    }
}
```

Lab: after midterm (Lab no:8)



Gives infix form of expression!

1-2-4

1-3

$1 - \binom{2}{2} - 4$   
 $1 - \binom{2}{2} - 5$

$\frac{9-1+1}{2}$

# EULER TOUR TRAVERSAL: GENERIC

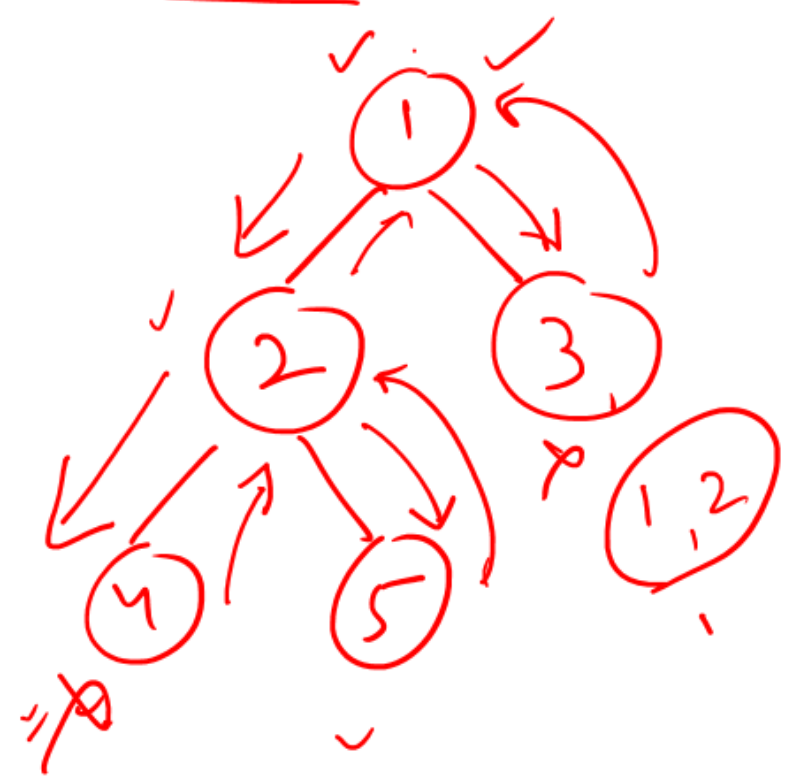
Applications: finding ancestor, LCA (lowest common ancestor), finding number of descendants, etc.

1	2	3	4	5	6	7	8	9
1	2	4	2	5	2	1	3	1

1	→	1	tin
2	→	2	tin
3	→	8	
4	→	3	
5	→	5	

tin(A) < tin(B)

twa(A) > twa(B)



# BINARY TREE CONSTRUCTION FROM TRAVERSAL ORDER

Can you construct the binary tree, given two traversal sequences?

preorder = ab ✓ //  
postorder = ba

Postorder  
index ←



do they uniquely define a binary tree?

inorder = g d h b e i a f j c ✓  
preorder = a b d g h e i c f j

Try:

Postorder: DEBFCA

Inorder: DBEAFC



# BINARY TREE ADT

*p.left()*: Return the left child of *p*; an error condition occurs if *p* is an external node.

*p.right()*: Return the right child of *p*; an error condition occurs if *p* is an external node.

*p.parent()*: Return the parent of *p*; an error occurs if *p* is the root.

*p.isRoot()*: Return true if *p* is the root and false otherwise.

*p.isExternal()*: Return true if *p* is external and false otherwise.

*size()*: Return the number of nodes in the tree.

*empty()*: Return true if the tree is empty and false otherwise.

*root()*: Return a position for the tree's root; an error occurs if the tree is empty.

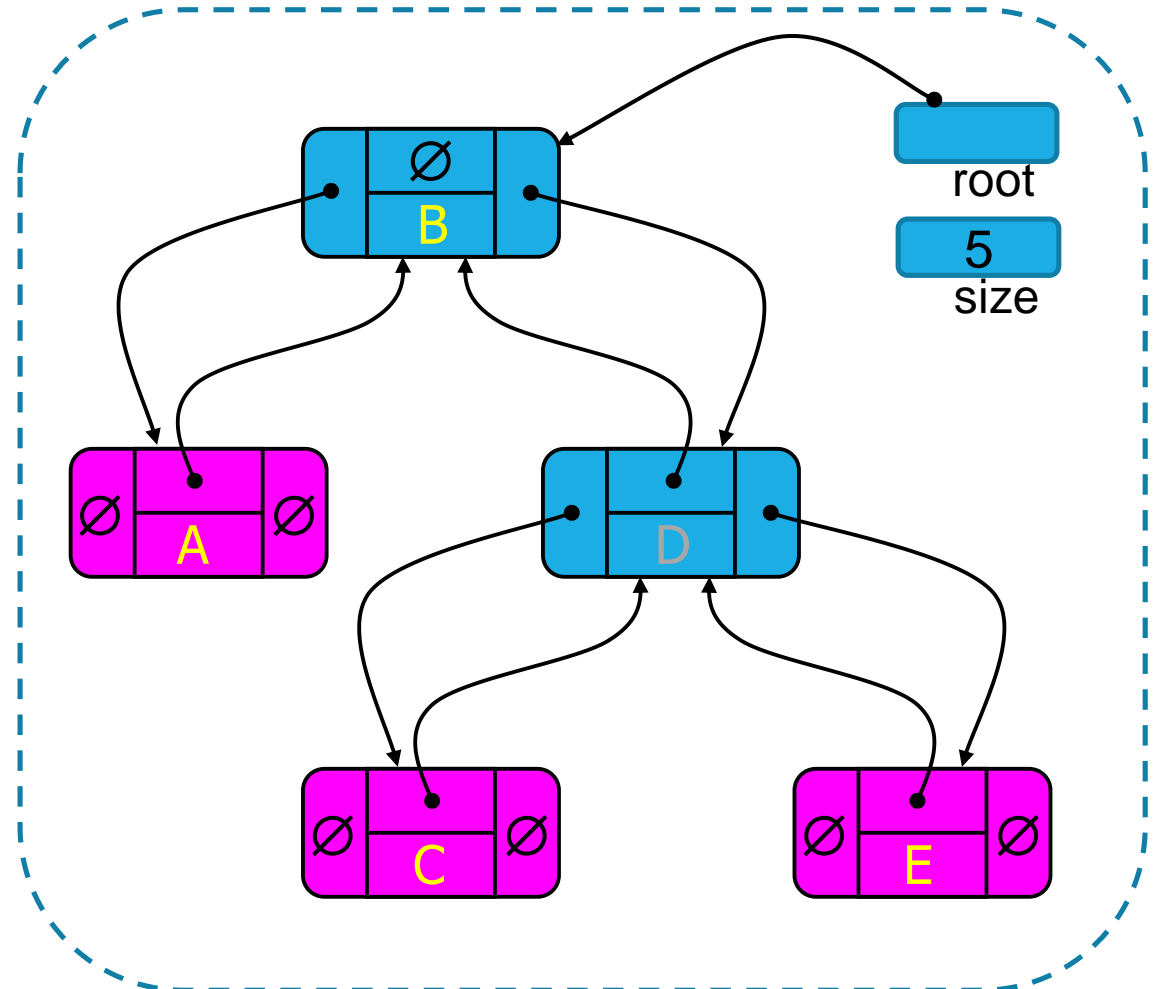
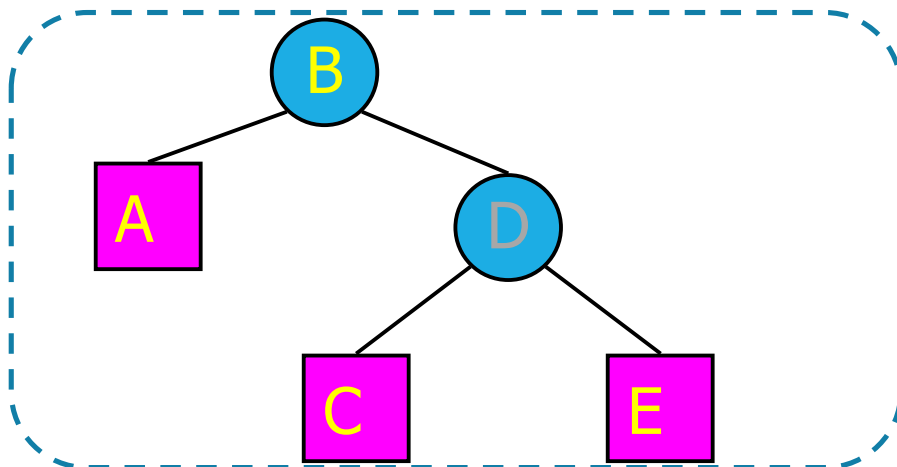
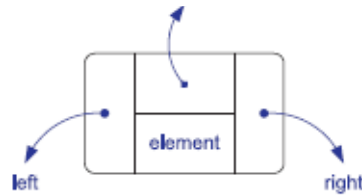
*positions()*: Return a position list of all the nodes of the tree.

```
template <typename E>
class Position<E> {
public:
    E& operator*();
    Position left() const;
    Position right() const;
    Position parent() const;
    bool isRoot() const;
    bool isExternal() const;
};
```

```
template <typename E>
class BinaryTree<E> {
public:
    class Position;
    class PositionList;
public:
    int size() const;
    bool empty() const;
    Position root() const;
    PositionList positions() const;
};
```

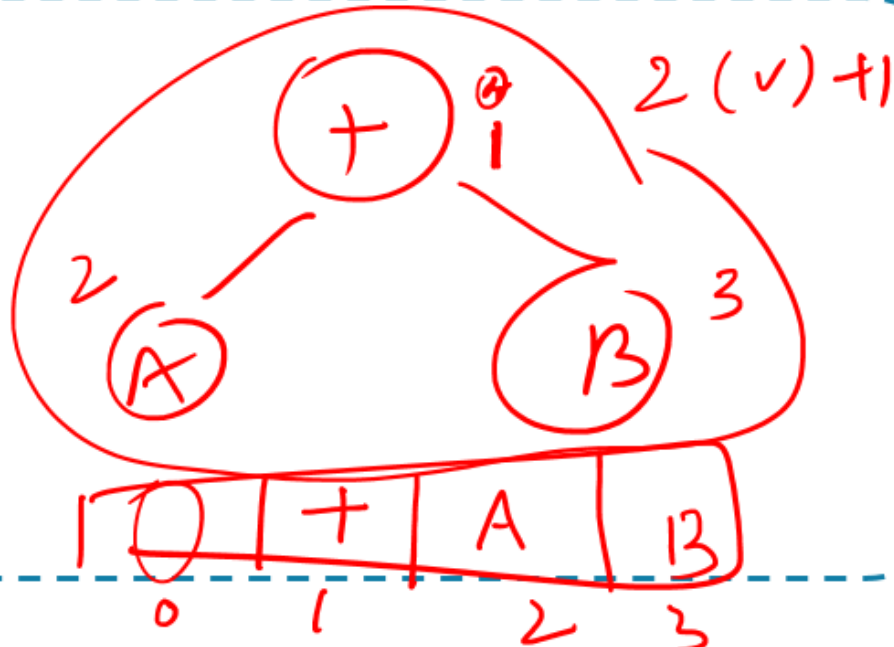
# LINKED STRUCTURE FOR BINARY TREES

```
struct Node{  
    Node *left;  
    Node *right;  
    Node *par;  
    int data;  
    Node() : data(), par(NULL), left(NULL), right(NULL) { }  
};
```

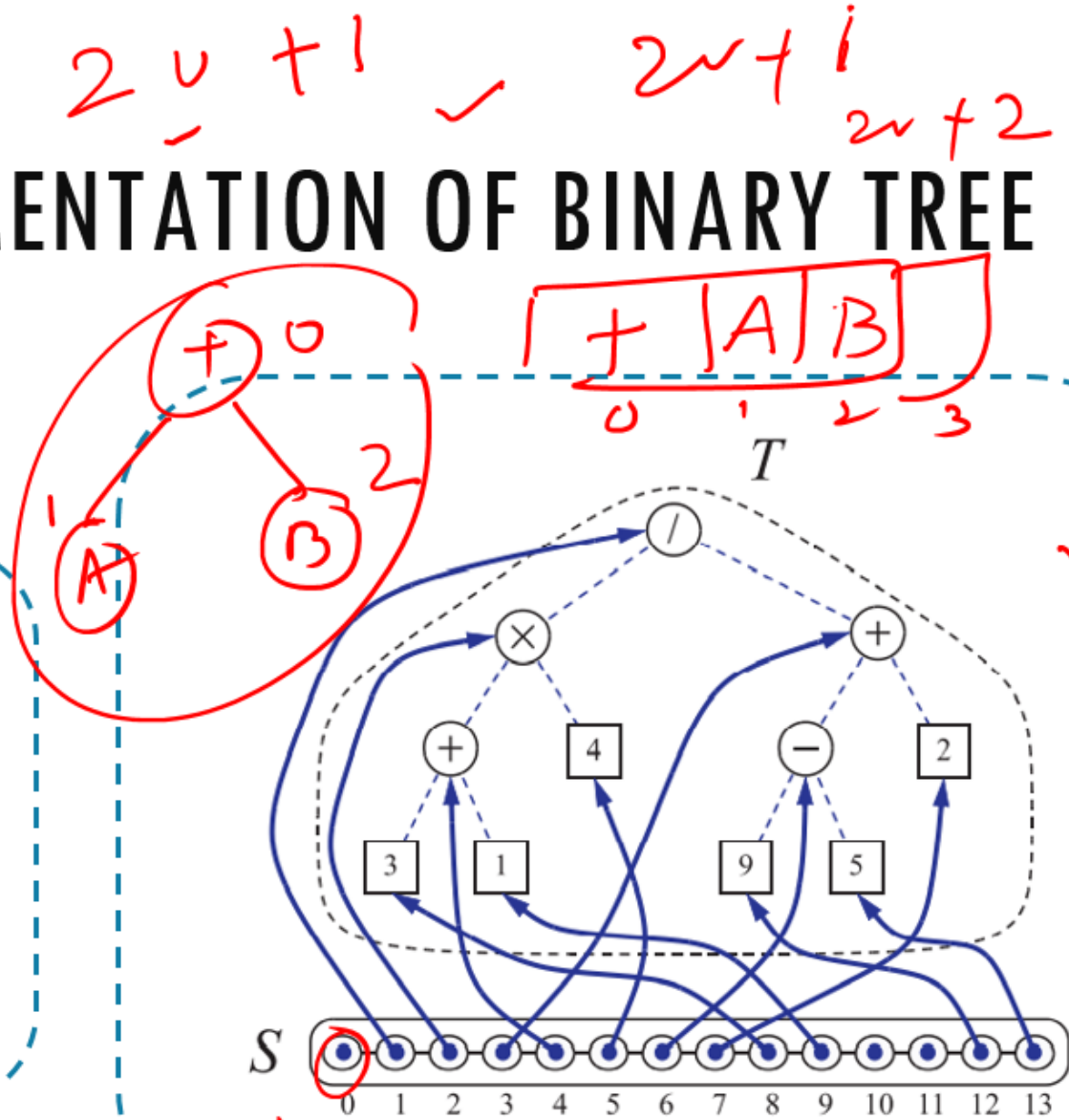


# A VECTOR-BASED IMPLEMENTATION OF BINARY TREE

- Level numbering scheme.



(An Example Expression tree)



(Example binary tree using a vector)





# BINARY TREE IMPLEMENTATION USING LINKED STRUCT

```

77 void BinaryTree::inorder(Node *ptr){
78     if(!ptr) return;
79     inorder(ptr->left);
80     cout<<" "<<ptr->data;
81     inorder(ptr->right);
82 }
83 void BinaryTree::postorder(Node *ptr){
84     if(!ptr) return;
85     postorder(ptr->left);
86     postorder(ptr->right);
87     cout<<" "<<ptr->data;
88 }
89 void BinaryTree::preorder(Node *ptr){
90     if(!ptr) return;
91     cout<<ptr->data<<" ";
92     preorder(ptr->left);
93     preorder(ptr->right);
94 }

```

```

98 Node* BinaryTree::createTree(vector<int> &v,Node *root,
99     Node *parent,int i){
100     n = v.size();
101     if(i<v.size()){
102         Node *temp = new Node;
103         temp->data = v[i];
104         temp->par = parent;
105         root = temp;
106         root->left = createTree(v,root->left,root,2*i+1);
107         root->right = createTree(v,root->right,root,2*i+2);
108         all_nodes.insert(root);
109     }
110     main_root = root;
111     return root;
112 }

```

```

Enter size of input array : 6
Enter array : 5 7 8 3 4 9
1
Size : 6
2
Tree is not empty
3
Inorder traversal : 3 7 4 5 9 8
4
Preorder traversal : 5 7 3 4 8 9
5
Postorder traversal : 3 4 7 9 8 5
6
Height by height1 : 2
7
Height by height2 : 2

```

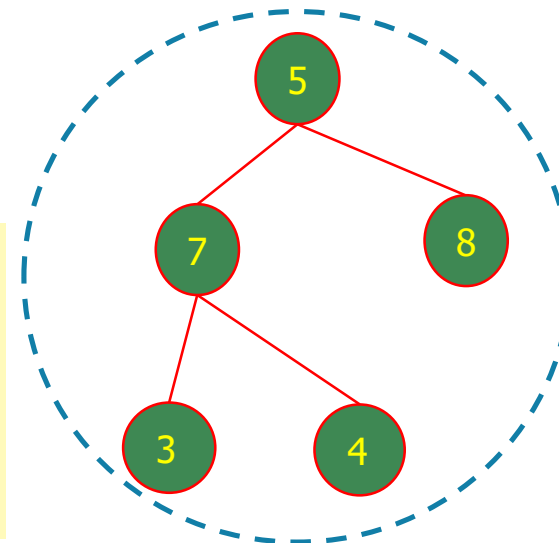
expandExternal(const Position& p)

## Lab 9

```

// expand external node
void LinkedBinaryTree::expandExternal(const Position& p) {
    Node* v = p.v;
    v->left = new Node;
    v->left->par = v;
    v->right = new Node;
    v->right->par = v;
    n += 2;
}
// p's node
// add a new left child
// v is its parent
// and a new right child
// v is its parent
// two more nodes

```



```

Enter size of input array : 5
Enter array : 5 7 8 3 4
1
Size : 5
2
Tree is not empty
3
Inorder traversal : 3 7 4 5 8
4
Preorder traversal : 5 7 3 4 8
5
Postorder traversal : 3 4 7 8 5
6
Height by height1 : 2
7
Height by height2 : 2

```

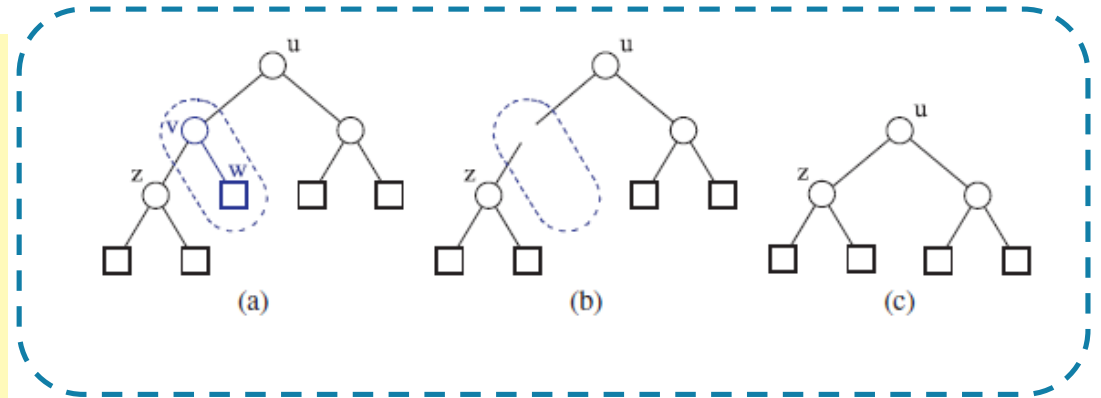
# BINARY TREE UPDATE FUNCTIONS

removeAboveExternal (const Position& p)

```

LinkedBinaryTree::Position // remove p and parent
LinkedBinaryTree::removeAboveExternal(const Position& p) {
    Node* w = p.v; Node* v = w->par; // get p's node and parent
    Node* sib = (w == v->left ? v->right : v->left); // child of root?
    if (v == _root) { // ...make sibling root
        _root = sib;
        sib->par = NULL;
    }
    else {
        Node* gpar = v->par; // w's grandparent
        if (v == gpar->left) gpar->left = sib; // replace parent by sib
        else gpar->right = sib;
        sib->par = gpar;
    }
    delete w; delete v; // delete removed nodes
    n -= 2; // two fewer nodes
    return Position(sib);
}

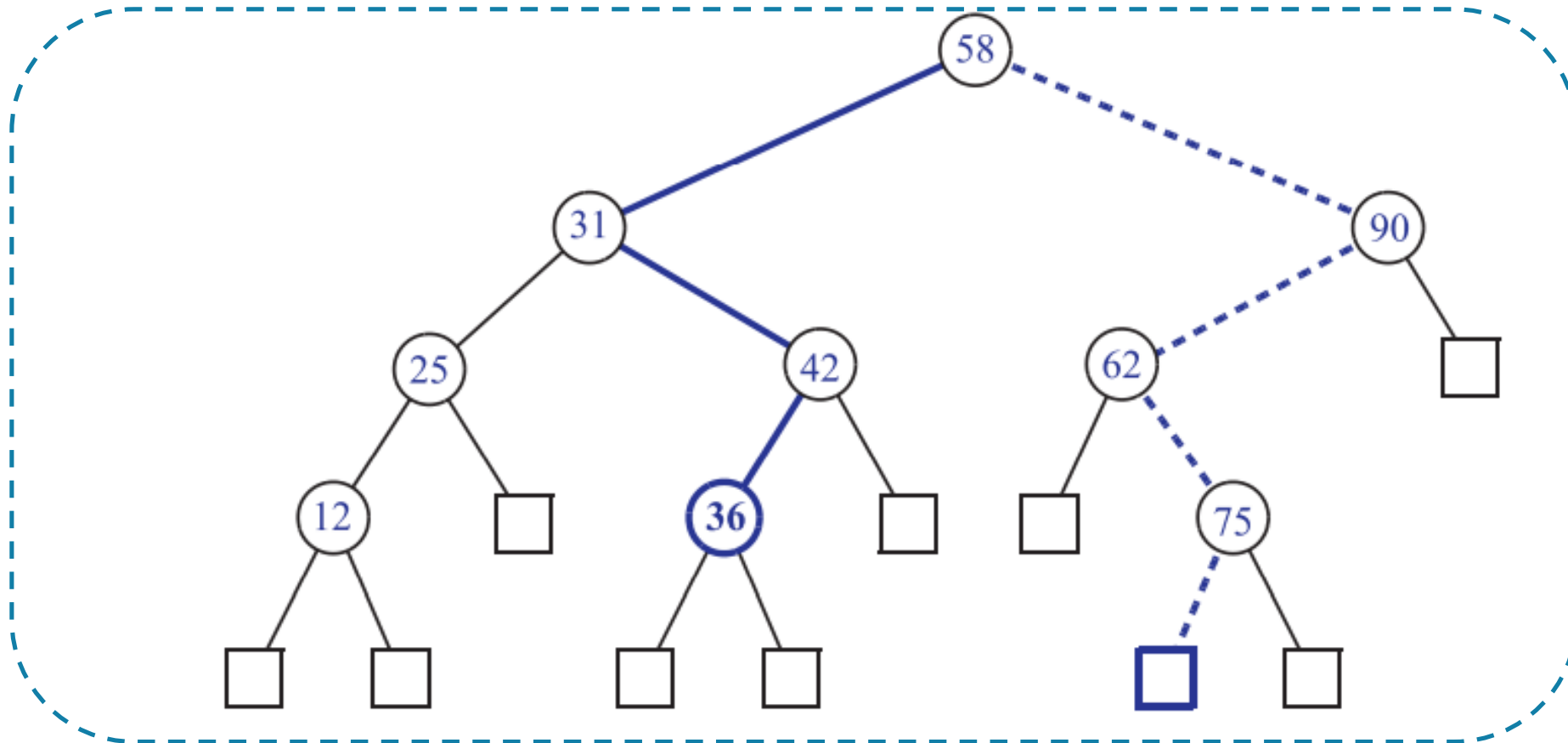
```



The vector implementation of a binary tree is a fast and easy way of realizing the binary-tree ADT, but it can be very space inefficient if the height of the tree is large. →  $O(2^n)$ , where 'n' is no. of nodes in T.

Operation	Time
left, right, parent, isExternal, isRoot	$O(1)$
size, empty	$O(1)$
root	$O(1)$
expandExternal, removeAboveExternal	$O(1)$
positions	$O(n)$

# BINARY SEARCH TREE (BST)



More later

...

(Searching successfully 36 through blue solid path, and unsuccessfully 71 through dashed path)

# THE TEMPLATE FUNCTION PATTERN

- The **template function pattern** describes a generic computation method that can be tuned for a particular application by redefining certain steps.

```
template <typename E, typename R>           // element and result types
class EulerTour {                          // a template for Euler tour
protected:
    struct Result {                         // stores tour results
        R leftResult;                      // result from left subtree
        R rightResult;                     // result from right subtree
        R finalResult;                     // combined result
    };
    typedef BinaryTree<E> BinaryTree;       // the tree
    typedef typename BinaryTree::Position Position; // a position in the tree
protected:                                // data member
    const BinaryTree* tree;                // pointer to the tree
public:
    void initialize(const BinaryTree& T)    // initialize
    { tree = &T; }
protected:                                // local utilities
    int eulerTour(const Position& p) const; // perform the Euler tour
                                        // functions given by subclasses
    virtual void visitExternal(const Position& p, Result& r) const {}
    virtual void visitLeft(const Position& p, Result& r) const {}
    virtual void visitBelow(const Position& p, Result& r) const {}
    virtual void visitRight(const Position& p, Result& r) const {}
    Result initResult() const { return Result(); }
    int result(const Result& r) const { return r.finalResult; }
};
```

(Class **EulerTour** defining a generic Euler tour of a binary tree. It realizes **template function pattern** and must be specialized for use)

```
template <typename E, typename R>           // do the tour
int EulerTour<E, R>::eulerTour(const Position& p) const {
    Result r = initResult();
    if (p.isExternal()) {                  // external node
        visitExternal(p, r);
    }
    else {                                  // internal node
        visitLeft(p, r);
        r.leftResult = eulerTour(p.left()); // recurse on left
        visitBelow(p, r);
        r.rightResult = eulerTour(p.right()); // recurse on right
        visitRight(p, r);
    }
    return result(r);
}
```

(Member function **eulerTour** recursively traverses the tree and accumulates the results)

# CONTINUED...

```
template <typename E, typename R>
class EvaluateExpressionTour : public EulerTour<E, R> {
protected: // shortcut type names
    typedef typename EulerTour<E, R>::BinaryTree BinaryTree;
    typedef typename EulerTour<E, R>::Position Position;
    typedef typename EulerTour<E, R>::Result Result;
public:
    void execute(const BinaryTree& T) { // execute the tour
        initialize(T);
        std::cout << "The value is: " << eulerTour(T.root()) << "\n";
    }
protected: // leaf: return value
    virtual void visitExternal(const Position& p, Result& r) const
    { r.finalResult = (*p).value(); }

    // internal: do operation
    virtual void visitRight(const Position& p, Result& r) const
    { r.finalResult = (*p).operation(r.leftResult, r.rightResult); }
};
```

(Implementation of class `EvaluateExpressionTour` which specializes `EulerTour` to evaluate the expression associated with an arithmetic expression tree)

```
template <typename E, typename R>
class PrintExpressionTour : public EulerTour<E, R> {
protected: // ...same type name shortcuts as in EvaluateExpressionTour
public:
    void execute(const BinaryTree& T) { // execute the tour
        initialize(T);
        cout << "Expression: "; eulerTour(T.root()); cout << endl;
    }
protected: // leaf: print value
    virtual void visitExternal(const Position& p, Result& r) const
    { (*p).print(); }

    // left: open new expression
    virtual void visitLeft(const Position& p, Result& r) const
    { cout << "("; }

    // below: print operator
    virtual void visitBelow(const Position& p, Result& r) const
    { (*p).print(); }

    // right: close expression
    virtual void visitRight(const Position& p, Result& r) const
    { cout << ")"; }
};
```

(A derived class, called `PrintExpressionTour` that prints the arithmetic expression)



# THANK YOU!

Next class: Priority Queues, and Heaps...